# M-Ary Data Transmission

#### Topics to be covered in this video and subsequent ones

- Orthogonal Functions and Signal Space Representation
- Gram-Schmidt Orthogonaliztion Procedure
- Optimum Receiver for Binary Transmission (Revisited using signal space concept)
- Optimum Receiver for M-Ary Transmission (using signal space representation)
- M-ary Coherent Amplitude-Shift Keying (M-ASK)
- M-ary Coherent Phase-Shift Keying (M-PSK)
- M-ary Coherent Frequency-Shift Keying (M-FSK)
- M-ary Quadrature Amplitude Modulation (M-QAM)
- Union Bound on the Symbol Probability of Error
- Comparison of the various M-ary modulation techniques

#### The Binary Communication System (Revisited)



# Assumptions

- In binary data transmission over a communication channel, logic 1 is represented by a signal  $s_1(t)$  and logic 0 by a signal  $s_2(t)$ .
- The time allocated for each signal is the bit duration  $T_b(\tau \text{ in the previous chapter})$  in the case of binary and  $T_s$  for the case of M-ary.
- The data rate is  $R_b = 1/T_b$  bits/sec.
- The channel noise n(t) is additive white Gaussian (AWGN) with a double-sided PSD of  $N_0/2$  W/Hz, mean  $E\{n(t)\} = 0$ ,  $R_n(\tau) = \frac{N_0}{2}\delta(\tau)$ . Noise is assumed to be added at the front end of the receiver (with variance  $N_0^2/2$ ).
- The data component at the front end of the receiver is assumed to be an exact replica of the transmitted signal, in the sense that the transmission bandwidth of the medium is wide enough to reproduce the signal without distortion.
- Bits in different time intervals are assumed independent.
- The signal to be processed by the receiver is the noisy signal  $y(t) = s_i(t) + n(t)$
- Based on y(t), the task of the receiver is to decide whether a 1 or a 0 was transmitted during each transmission slot  $\tau$  with minimum probability of error.
- The approach is based on the signal space where time functions are represented by numbers, which may be deterministic or random depending on the type of signal.

#### Geometric Representation of Signals (Signal Space Concept)



Signal Space Representation: Energy, Distance, and Probability of Error

• 
$$E_1 = \int_0^{T_b} (s_1(t)^2 dt) = \int_0^{T_b} (s_{11} \phi_1(t) + s_{12} \phi_2(t))^2 dt$$
,  
•  $E_1 = \int_0^{T_b} (s_1(t)^2 dt) = (s_{11})^2 + (s_{12})^2$ , Signal Energy  
•  $E_2 = \int_0^{T_b} (s_2(t)^2 dt) = \int_0^{T_b} (s_{21} \phi_1(t) + s_{22} \phi_2(t))^2 dt$   
•  $E_2 = \int_0^{T_b} (s_2(t)^2 dt) = (s_{21})^2 + (s_{22})^2$ , Signal Energy  
•  $d_{12}^2 = \int_0^{T_b} (s_1(t) - s_2(t))^2 dt = (s_{11} - s_{21})^2 + (s_{12} - s_{22})^2$ , Square of the Distance Between two Signals

• 
$$P_b^* = Q\left(\sqrt{\frac{\int_0^{T_b} (s_1(t) - s_2(t))^2 dt}{2N_0}}\right) = Q\left(\frac{d_{12}}{\sqrt{2N_0}}\right);$$
 Bit Error Probability

## Geometric Representation of Signals: Summary $\phi_2(t)$ $s_1(t) = s_{11}\phi_1(t) + s_{12}\phi_2(t),$ **Signal Representation** $(S_{11}, S_{12})$ $s_2(t) = s_{21}\phi_1(t) + s_{22}\phi_2(t),$ $s_{ij} = \int_{0}^{T_b} s_i(t)\phi_j(t) dt, \ i, j \in \{1, 2\}, \$ Signal Coefficients $d_{12}^{2} = \int_{0}^{T_{b}} (s_{1}(t) - s_{2}(t))^{2} dt$ $= (s_{11} - s_{21})^{2} + (s_{12} - s_{22})^{2}$ $\rightarrow \phi_{\rm l}(t)$ 0 *S*<sub>11</sub> S 21 $E_{1} = \int_{0}^{T_{b}} (s_{1}(t)^{2} dt = (s_{11})^{2} + (s_{12})^{2}$ $E_{2} = \int_{0}^{T_{b}} (s_{2}(t)^{2} dt = (s_{21})^{2} + (s_{22})^{2}$ $P_{b}^{*} = Q\left(\sqrt{\frac{\int_{0}^{T_{b}} (s_{1}(t) - s_{2}(t))^{2} dt}{2N_{0}}}\right) = Q\left(\frac{d_{12}}{\sqrt{2N_{0}}}\right)$

**Gram-Schmidt Method:** Basis for a Two-Dimensional Space

- The Gram-Schmidt method is a procedure for generating a set of orthonormal basis functions from a given set of functions  $(s_1(t), s_2(t))$
- The original set of functions may be dependent or independent, but the basis functions are both
- Linearly independent and
- Orthonormal.



#### Gram-Schmidt Method: Basis for a Two-Dimensional Space



# Gram-Schmidt Method: M-Ary Case $\phi_1(t) = \frac{s_1(t)}{\sqrt{\int_{-\infty}^{\infty} s_1^2(t) \mathrm{d}t}},$ $\phi_i(t) = \frac{\phi'_i(t)}{\sqrt{\int_{-\infty}^{\infty} \left[\phi'_i(t)\right]^2 \mathrm{d}t}}, \quad i = 2, 3, \dots, N,$ $\phi'_{i}(t) = \frac{s_{i}(t)}{\sqrt{E_{i}}} - \sum_{i=1}^{i-1} \rho_{ij}\phi_{j}(t),$ $\rho_{ij} = \int_{-\infty}^{\infty} \frac{s_i(t)}{\sqrt{E_i}} \phi_j(t) \mathrm{d}t, \quad j = 1, 2, \dots, i-1.$

If the waveforms  $\{s_i(t)\}_{i=1}^M$  form a *linearly independent set*, then N = M. Otherwise N < M.

#### **Example: Polar Non-return to zero Binary Signals**



 $s_2(t) = -s_1(t)$ ; Linearly Dependent

$$\rho = \frac{1}{\sqrt{E_1 E_2}} \int_0^{T_b} s_1(t) s_2(t) dt = -1$$

For the case of binary antipodal signaling, i.e., when  $s_2(t) = -s_1(t)$ , we need only one basis function since signals are linearly dependent. The signals are represented as:  $s_1(t) = \sqrt{E}\phi_1(t)$  $s_2(t) = -\sqrt{E}\phi_1(t)$ 

(a) Signal set. (b) Orthonormal function. (c) Signal space representation.

**Example: Orthogonal Binary Signals** 



#### Signal Space Representation of Signals: Revisited

 $\phi_2(t)$ 

S11 S21

In the previous video, we saw that if a signal space is characterized by two basis functions  $\emptyset_1(t)$  and  $\emptyset_2(t)$ , and if we are given two signals  $s_1(t)$  and  $s_2(t)$ , then  $s_1(t)$  and  $s_2(t)$  can be expressed as:

$$s_{1}(t) = s_{11}\phi_{1}(t) + s_{12}\phi_{2}(t), \qquad (1)$$

$$s_{2}(t) = s_{21}\phi_{1}(t) + s_{22}\phi_{2}(t).$$
where  $s_{ij} = \int_{0}^{T_{b}} s_{i}(t)\phi_{j}(t)dt, \quad i, j \in \{1, 2\}, \quad (2)$ 

$$s_{11}, s_{12}; \qquad (s_{11}, s_{12}); \qquad (s_{21}, s_{22}); \qquad (s_{21}, s_{22$$

- If the receiver is able to recover the coefficients  $s_{ij}$  using (2), then according to (1) the signals are completely known. This is of course true in the absence of noise. However, in the presence of noise, a random variable is added to  $s_{ij}$  that makes the decision more involved.
- The objective of the receiver is to retrieve the coefficients **s**<sub>ii</sub> using (2) and to make a decision accordingly

#### **Receiver Structure**

- Since each signal is characterized by two coefficients in the  $\emptyset_1(t)$  and  $\emptyset_2(t)$  plane, then we need two correlators that retrieve these coefficients at the receiver (figure below)
- However, in the presence of noise, the received signal is the sum of the transmitted signal and the AWGN component. That is:  $r(t) = s_i(t) + w(t)$ .
- The components in the  $\phi_1(t)$  and  $\phi_2(t)$  directions are:

 $r_{1} = \int_{0}^{T_{b}} (s_{i}(t) + w(t)) \phi_{1}(t) dt = s_{i1} + N_{1}; N_{1} \sim N(0, N_{0}|2); r_{1} \sim N(s_{i1}, N_{0}|2); r_{1} and r_{2} are$  $r_{2} = \int_{0}^{T_{b}} (s_{i}(t) + w(t)) \phi_{2}(t) dt = s_{i2} + N_{2}; N_{2} \sim N(0, N_{0}|2); r_{2} \sim N(s_{i2}, N_{0}|2); independent$ 



#### The Minimum Probability of Error Receiver

• Given the correlator outputs  $r_1$  and  $r_2$ , we need to find the decision rule that minimizes the probability of error

$$r_{1} = \int_{0}^{T_{b}} (s_{i}(t) + w(t)) \phi_{1}(t) dt = s_{i1} + N_{1}; N_{1} \sim N(0, N_{0}|2); r_{1} \sim N(s_{i1}, N_{0}|2); r_{1} and r_{2} are$$
  

$$r_{2} = \int_{0}^{T_{b}} (s_{i}(t) + w(t)) \phi_{2}(t) dt = s_{i2} + N_{2}; N_{2} \sim N(0, N_{0}|2); r_{2} \sim N(s_{i2}, N_{0}|2); independent$$

• The decision rule will be function of the received observations  $r_1$  and  $r_2$ , the signal components  $s_{ij}$ , and the noise power  $N_0$  as we shall derive next.



#### The Optimum Decision Rule

- Let P<sub>1</sub> and P<sub>2</sub> be the probability of sending signals s<sub>1</sub> and s<sub>2</sub>, respectively.
- Let R<sub>1</sub> and R<sub>2</sub> be the decision regions in the two-dimensional space corresponding to signals s<sub>1</sub> and s<sub>2</sub>, respectively.
- If  $(r_1, r_2 \in R_1)$  decide  $s_1$  (digit 1) Also. Else, decide  $s_2$  (digit 0).

 $P_b = P(send s_1, decide s_2) + P(send s_2, decide s_1)$ 

 $P_b = P_1 P(decide s_2|s_1) + P_2 P(decide s_1|s_2)$ 



#### The Optimum Decision Rule

$$P_{b} = P_{1} \int_{R_{2}} f(r_{1}, r_{2} | b_{i} = 1) dr_{1} dr_{2} + P_{2} \int_{R_{1}} f(r_{1}, r_{2} | b_{i} = 0) dr_{1} dr_{2}$$

$$P_{b} = P_{1} \int_{R_{2}} f(r_{1}, r_{2} | b_{i} = 1) dr_{1} dr_{2} + P_{2} (\int_{R} f(r_{1}, r_{2} | b_{i} = 0) dr_{1} dr_{2} - \int_{R_{2}} f(r_{1}, r_{2} | b_{i} = 0) dr_{1} dr_{2})$$

$$P_{b} = P_{2} \int_{R} f(r_{1}, r_{2} | b_{i} = 0) dr_{1} dr_{2}$$

$$+ \int_{R_{2}} \{P_{1}f(r_{1}, r_{2} | b_{i} = 1) - P_{2}f(r_{1}, r_{2} | b_{i} = 0)\} dr_{1} dr_{2}$$

$$(r_{1}, r_{2})$$

$$R_{2} \int_{(s_{21}, s_{22})}^{s_{2}} (r_{1}, r_{2} | b_{i} = 1) - P_{2}f(r_{1}, r_{2} | b_{i} = 1) < P_{2}f(r_{1}, r_{2} | b_{i} = 0)$$

$$R_{2} \text{ assign to } R_{2} \text{ all values of } (r_{1}, r_{2} | b_{i} = 1) < P_{2}f(r_{1}, r_{2} | b_{i} = 0)$$

$$Decide 0 \text{ when } P_{1}f(r_{1}, r_{2} | b_{i} = 1) < P_{2}f(r_{1}, r_{2} | b_{i} = 0)$$

$$\frac{f(r_{1}, r_{2} | b_{i} = 1)}{f(r_{1}, r_{2} | b_{i} = 0)} > \frac{P_{2}}{P_{1}}, Decide 1, otherwise decide 0$$
This is known as the likelihood ratio test

#### **Receiver Structure**

- The probability of error is minimized when the following decision rule is employed:
- Decide 1 when  $\frac{f(r_1, r_2 | b_i = 1)}{f(r_1, r_2 | b_i = 0)} \ge \frac{P_2}{P_1}$ ; (1); else, decide 0;
- $f(r_1, r_2|b_i = 1) = f(r_1|b_i = 1)f(r_2|b_i = 1)$ ; due to independence ;  $X \sim N(\mu, \sigma^2 = N_0 |2)$ ;  $\frac{1}{\sqrt{2\pi\sigma^2}} e^{\frac{(x-\mu)^2}{2\pi\sigma^2}}$
- $f(r_1, r_2|b_i = 0) = f(r_1 | b_i = 0)f(r_2|b_i = 0)$ ; due to independence
- $f(r_1 | b_i = 1) \sim N\left(s_{11}, \frac{N_0}{2}\right); ; f(r_2 | b_i = 1) \sim N\left(s_{12}, \frac{N_0}{2}\right);$
- $f(r_1 | b_i = 0) \sim N\left(s_{21}, \frac{N_0}{2}\right); ; f(r_2 | b_i = 0) \sim N\left(s_{22}, \frac{N_0}{2}\right);$



#### **Optimum Receiver: Binary Case**

$$(r_1 - s_{11})^2 + (r_2 - s_{12})^2 \stackrel{\stackrel{1_D}{\geq}}{\underset{0_D}{\geq}} (r_1 - s_{21})^2 + (r_2 - s_{22})^2 + N_0 \ln\left(\frac{P_1}{P_2}\right)$$

• For the special case of  $P_1 = P_2$  (signals are equally likely):

$$(r_1 - s_{11})^2 + (r_2 - s_{12})^2 \stackrel{1_D}{\underset{0_D}{\geq}} (r_1 - s_{21})^2 + (r_2 - s_{22})^2.$$

## This is the minimum distance decision rule

The probability of error for Equally-probable signals is given by:

$$P_b^* = Q\left(\frac{d_{12}}{\sqrt{2N_0}}\right) = Q\left(\sqrt{\frac{\int_0^{T_b} (s_1(t) - s_2(t))^2 dt}{2N_0}}\right)$$
$$d_{12}^2 = \int_0^{T_b} (s_1(t) - s_2(t))^2 dt = (s_{11} - s_{21})^2 + (s_{12} - s_{22})^2$$

As derived earlier

**Optimum Receiver : Matched Filter and Correlators** 





The receiver can be implemented in terms of correlators and can, as well, be implemented in terms of the matched filters. Here, matched means that the filters at the receiver are matched to the basis functions used in the transmission process. The two figures on this slide are equivalent in terms of performance. 8

#### Summary of Results on the Binary Case



#### M-Ary Transmission

- In M-ary transmission, a block of n binary digits are grouped together to form one symbol (message).
- If  $T_b$  is the bit duration, then  $T_s = nT_b$  is the symbol duration. The data rates are related by:  $R_s = R_b/n$ .
- There are  $M = 2^n$  possible symbols. Hence, we need  $M = 2^n$  signals to be transmitted.
- The signals can modulate a high frequency carrier in the amplitude, the phase, frequency, and both the amplitude and the phase.
- We will study the following modulation techniques:
- M-ary ASK, M-ary PSK, M-ary FSK, and Quadrature Amplitude Modulation (QAM).
- For each modulation scheme, we will consider the transmitter, the optimum receiver, the probability of error, the power spectral density and the bandwidth.

#### M-Ary Transmission

- In most of our analysis here, we will encounter M-ary transmission in a twodimensional space (except for the M-ary FSK), in which we need two basis functions  $\emptyset_1(t)$  and  $\emptyset_2(t)$ .
- In this space, the signals are represented as:
- $s_1(t) = s_{11} \phi_1(t) + s_{12} \phi_2(t)$ •  $s_3(t) = s_{31} \phi_1(t) + s_{32} \phi_2(t)$ • Where  $s_{i1} = \int_0^{T_s} s_i(t) \phi_1(t) dt$ ,  $s_{i2} = \int_0^{T_s} s_i(t) \phi_2(t) dt$
- The receiver has to decide on which signal was transmitted based on the received vector  $(r_1, r_2)$ .
- Note: To obtain the basis functions from the M given functions, one can use the Gram Schmidt orthogonalization procedure. The number of basis functions N<=M.</li>

#### Optimum Receiver for M-Ary Transmission

- The observation space is to be partitioned into M regions, such that if the set of measurements fall into region  $R_k$  signal  $s_k$  is declared true.
- It is assumed here that all signals are equally probable.
- The receiver collects the measurements from the N correlators (r vector) and calculates the distance to each of the N signals.
- It decides in favor of the signal closest to the (r vector).



#### M-ary Coherent Amplitude-Shift Keying (M-ASK)

$$s_{i}(t) = V_{i}\sqrt{\frac{2}{T_{s}}}\cos(2\pi f_{c}t), \ 0 \leq t \leq T_{s} \quad s_{i}(t) = V_{i}\phi_{1}(t)$$

$$= (i-1)\Delta$$

$$= [(i-1)\Delta]\phi_{1}(t), \ \phi_{1}(t) = \sqrt{\frac{2}{T_{s}}}\cos(2\pi f_{c}t), \ 0 \leq t \leq T$$

$$i = 1, 2, \dots, M. \quad E_{i} = ((i-1)\Delta)^{2} \quad E_{i} = (V_{i})^{2}$$

$$\xrightarrow{s_{1}(t)} \underbrace{s_{2}(t)}_{\Delta} \underbrace{s_{3}(t)}_{2\Delta} \dots \underbrace{s_{k}(t)}_{(k-1)\Delta} \underbrace{m_{i-1}(t)}_{(M-2)\Delta} \underbrace{m_{i}(t)}_{(M-1)\Delta} \phi_{1}(t)$$

$$\xrightarrow{s_{i}(t)} \underbrace{r(t)}_{\psi_{1}(t)} \underbrace{f_{i}(t)}_{(k-1)T_{i}} \underbrace{f_{i}(t)}_{E_{2}} \underbrace{f_{i}(t)}_{E_{2}} \underbrace{m_{i}(t)}_{T_{s}} \underbrace{f_{i}(t)}_{E_{2}} \underbrace{m_{i}(t)}_{T_{s}} \underbrace{f_{i}(t)}_{E_{2}} \underbrace{f_{i}(t)}_{T_{s}} \underbrace{f_{i}(t)}_{E_{s}(t)} \underbrace{f_{i}(t)}_{E_$$

In this case, we have M signals. However, we need only one base s, function. The signals are linearly dependent and hence, every signal can be expressed in terms of this base function.

Since there is one base function, the receiver consists of one correlator (multiplier followed by an integrator), a sampler, and a decision device (set of comparators).

#### Minimum-Distance Decision Rule for M-ASK

 $\mathsf{Choose} \left\{ \begin{array}{ll} s_k(t), & \text{if} \quad \left(k - \frac{3}{2}\right) \Delta < r_1 < \left(k - \frac{1}{2}\right) \Delta, \ k = 2, 3, \dots, M - 1 \\ s_1(t), & \text{if} \quad r_1 < \frac{\Delta}{2} \\ s_M(t), & \text{if} \quad r_1 > \left(M - \frac{3}{2}\right) \Delta \end{array} \right.$ 



#### Minimum Distance Rule and Error Probability for two signals



#### Minimum-Distance Decision and Error Probability for M-ASK $f(\mathbf{r}_1|\mathbf{s}_k(t))$ . . . $r_1$ $(k-1)\Delta$ 0 Δ Choose $s_1(t)$ $\Rightarrow$ Choose $s_{\mu}(t)$ Choose $s_{i}(t)$ $P[\text{error}] = \sum_{i=1}^{M} P[s_i(t)]P[\text{error}|s_i(t)]$ $P[\text{error}|s_i(t)] = 2Q\left(\Delta/\sqrt{2N_0}\right), \quad i = 2, 3, \dots, M-1$ $P[\operatorname{error}|s_i(t)] = Q\left(\Delta/\sqrt{2N_0}\right), \quad i = 1, M$ $P[\text{error}] = \frac{2(M-1)}{M} Q\left(\Delta/\sqrt{2N_0}\right).$

For a given M, P[error] depends on the noise power  $(N_0)$  and the *minimum* distance  $\delta$ . This means that moving the origin of the signal constellation does not affect the performance  $\geq$ 

#### **Modified M-ASK Constellation**

The maximum and average transmitted energies can be reduced, without any sacrifice in error probability, by changing the signal set to one which includes the negative version of each signal.  $E_i = (V_i)^2$ 

$$s_i(t) = \underbrace{(2i-1-M)\frac{\Delta}{2}}_{V_i} \sqrt{\frac{2}{T_s}} \cos(2\pi f_c t), \ 0 \le t \le T_s, \ i = 1, 2, \dots, M.$$

(a) 
$$-\frac{3\Delta}{2}$$
  $-\frac{\Delta}{2}$   $0$   $\frac{\Delta}{2}$   $\frac{3\Delta}{2}$   $\cdots$ 

$$\begin{split} E_s &= \frac{\sum_{i=1}^M E_i}{M} = \frac{\Delta^2}{4M} \sum_{i=1}^M (2i-1-M)^2 = \frac{(M^2-1)\Delta^2}{12}.\\ E_b &= \frac{E_s}{\log_2 M} = \frac{(M^2-1)\Delta^2}{12\log_2 M} \Rightarrow \Delta = \sqrt{\frac{(12\log_2 M)E_b}{M^2 - 1}} \\ \end{split}$$

Es: Average Energy per Symbol

**Eb: Average Energy per bit** 

#### Probability of Symbol Error for M-ASK

$$P[\text{error}] = \frac{2(M-1)}{M} Q\left(\sqrt{\frac{6E_s}{(M^2-1)N_0}}\right) = \frac{2(M-1)}{M} Q\left(\sqrt{\frac{6\log_2 M}{M^2-1}\frac{E_b}{N_0}}\right).$$

$$P[\text{bit error}] = \frac{1}{\lambda} P[\text{symbol error}] = \frac{2(M-1)}{M\log_2 M} Q\left(\sqrt{\frac{6\log_2 M}{M^2-1}\frac{E_b}{N_0}}\right) \text{ (with Gray mapping)}$$

$$Bit error$$

$$P[\text{bit error}] = \frac{1}{\lambda} P[\text{symbol error}] = \frac{2(M-1)}{M\log_2 M} Q\left(\sqrt{\frac{6\log_2 M}{M^2-1}\frac{E_b}{N_0}}\right) \text{ (with Gray mapping)}$$

$$P[\text{bit error}] = \frac{1}{\lambda} P[\text{symbol error}] = \frac{2(M-1)}{M\log_2 M} Q\left(\sqrt{\frac{6\log_2 M}{M^2-1}\frac{E_b}{N_0}}\right) \text{ (with Gray mapping)}$$



Two comments: Error probability: for a given Eb/NO, increasing M results in an increase in the error probability.

Bandwidth: Increasing M results in a reduction in the bandwidth by a factor of  $\lambda = log_2(M)$ .

W is obtained by using the  $WT_s = 1$  rule-of-thumb. Here  $1/T_b$  is the bit rate (bits/s).

#### Example of 2-ASK (BPSK) and 4-ASK Signals

Baseband information signal 0 -1Tb 2Tb 3Tb 4Tb 5Tb 6Tb 7Tb 8Tb 9Tb 10Tb 0 BPSK Signalling 2 -2Tb 2Tb 3Tb 4Tb 5Tb 6Tb 7Tb 8Tb 9Tb 10Tb 0 4-ASK Signalling 2 -22Tb 0 4Tb 6Tb 8Tb 10Tb

Binary sequence: 1101101100

 $1 \rightarrow cos(2\pi f_0) t$  $0 \rightarrow -cos(2\pi f_0) t$ (similar to BPSK)

 $11 \rightarrow \cos(2\pi f_0) t$   $01 \rightarrow -\cos(2\pi f_0) t$   $10 \rightarrow 2\cos(2\pi f_0) t$  $00 \rightarrow -2\cos(2\pi f_0) t$ 

$$\begin{aligned} \text{M-ary Phase-Shift Keying (M-PSK)} \\ s_i(t) &= V \cos \left[ 2\pi f_c t - \frac{(i-1)2\pi}{M} \right], \quad 0 \leq t \leq T_s, \\ i &= 1, 2, \dots, M; \ f_c &= k/T_s, \ k \text{ integer; } E_s &= V^2 T_s/2 \text{ joules} \\ s_i(t) &= V \cos \left[ \frac{(i-1)2\pi}{M} \right] \cos(2\pi f_c t) + V \sin \left[ \frac{(i-1)2\pi}{M} \right] \sin(2\pi f_c t). \\ s_i(t) &= \frac{V \cos(2\pi f_c t)}{\sqrt{E_s}}, \ \phi_2(t) &= \frac{V \sin(2\pi f_c t)}{\sqrt{E_s}}. \\ s_{i1} &= \sqrt{E_s} \cos \left[ \frac{(i-1)2\pi}{M} \right], \ s_{i2} &= \sqrt{E_s} \sin \left[ \frac{(i-1)2\pi}{M} \right]. \\ \text{The signals lie on a circle of radius } \sqrt{E_s}, \text{ and are spaced every} \\ 2\pi/M \text{ radians around the circle.} \end{aligned}$$

# M-ary Phase-Shift Keying (M-PSK) $s_i(t) = V \cos\left[2\pi f_c t - \frac{(i-1)2\pi}{M}\right], \quad 0 \le t \le T_s, \quad \bullet \quad \text{Here, the amplitude of the carrier remains constant,}$

 $i = 1, 2, \ldots, M; f_c = k/T_s, k \text{ integer}; E_s = V^2 T_s/2 \text{ joules}$ 



- Here, the amplitude of the carrier remains constant, however the phase takes on one of M possible values.
- Two base functions are needed to represent all signals in the twodimensional signal space.
- The spacing between adjacent signals is  $\Delta \theta = 2\pi/M$  radians.
- In this example, M=8 and  $\Delta \theta = \frac{\pi}{4} = 45 \ degrees.$
- To minimize error, gray coding is used.

# M-ary Phase-Shift Keying (M-PSK): Signal Space Representation $s_i(t) = V \cos \left[ 2\pi f_c t - \frac{(i-1)2\pi}{M} \right], \quad 0 \le t \le T_s,$ $i = 1, 2, \dots, M; \ f_c = k/T_s, \ k \text{ integer}; \ E_s = V^2 T_s/2 \text{ joules}$



#### **Optimum receiver in a two-dimensional space**



 $r_1 \sim N(s_{i1}, N_0 | 2)$ ; Gaussian with mean  $s_{i1}$ , variance  $N_0 | 2$  $r_2 \sim N(s_{i2}, N_0 | 2)$ ; Gaussian with mean  $s_{i2}$ , variance  $N_0 | 2$ 

 $r_1 and r_2$  are independent

 $\begin{array}{l} \textit{Minimum Distance Rule} \\ \textit{Calculate: } d_1^2 = (r_1 - s_{11})^2 + (r_2 - s_{12})^2 \\ \textit{Calculate: } d_2^2 = (r_1 - s_{21})^2 + (r_2 - s_{22})^2 \\ \textit{Choose } s_1 \textit{ if } d_1^2 < d_2^2 \end{array}$ 

#### M-ary Phase-Shift Keying: Error Probability

 $\Pr[\text{error}] = \Pr[\text{error}|s_1(t)] = \Pr[r_1, r_2 \text{ fall outside Region } 1|s_1(t) \text{ transmitted}]$ 

$$= 1 - \Pr[r_1, r_2 \text{ fall in Region } 1|s_1(t) \text{ transmitted}]$$
$$= 1 - \iint_{r_1, r_2 \in \mathcal{R}_1} f(r_1, r_2|s_1(t)) dr_1 dr_2$$



$$d_{\min} = 2\sqrt{E_s} \sin\left(\frac{\pi}{M}\right)$$



#### Probability of Error in M-PSK

Region 2

Choose  $s_2(t)$ 

 $r_1$ 

Region 1

Choose  $s_1(t)$ 

 $s_2(t)$ 

 $\pi/M$ 

 $s_1(t)$ 

 $\sqrt{E}$ 

• The distance between two neighboring symbols is

$$\boldsymbol{P}_{\min} = 2\sqrt{E_s} \sin\left(\frac{\pi}{M}\right) \qquad \boldsymbol{P}_{\boldsymbol{b}}^* = \boldsymbol{Q}\left(\frac{\boldsymbol{d}_{12}}{\sqrt{2N_0}}\right)$$

• Each symbol has 2 close neighbor symbols.

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• An approximation for the symbol error prob.

 $P_s \approx (\text{Number of Signals at distance dmin}) Q \left(\frac{d_{\min}}{\sqrt{2N}}\right) = 2 \cdot Q \left(\frac{d_{\min}}{\sqrt{2}}\right)$ 

$$=2\cdot Q\left(\frac{2\sqrt{E_s}\sin\left(\frac{\pi}{M}\right)}{\sqrt{2N_0}}\right)=2\cdot Q\left(\sqrt{2\frac{E_s}{N_0}}\sin^2\left(\frac{\pi}{M}\right)}\right)$$

#### Probability of Error in M-PSK



#### Symbol and Bit Error Probability of M-PSK

- When Gray coding is used, the symbol and bit error probabilities are related by:  $P_b = \frac{1}{log_2(M)} P_s$ ;
- Moreover, the symbol energy is related to the bit energy by

• 
$$E_b = \frac{1}{\log_2(M)} E_S$$

- The performance of digital communication systems is usually taken as the error probability versus  $\frac{E_b}{N_0}$ .
- The next figure depicts the symbol probability of error for M-PSK

#### **Performance of M-PSK**



$$P_s \approx 2 \cdot Q \left( \sqrt{2 \frac{E_s}{N_0}} \sin^2 \left( \frac{\pi}{M} \right) \right)$$

As M increases, the symbol probability of error increases. Note that as M increases, the spacing between signals around the perimeter of the unit circle becomes smaller, and this results in a higher probability of error

#### M-ary Coherent Frequency-Shift Keying (M-FSK)

Signal Set:

$$s_m(t) = \sqrt{\frac{2E_s}{T_s}}\cos(2\pi(f_c + m\Delta f)t); m = 1, 2, ..., M, 0 \le t \le T_s$$

#### **Orthogonality condition:**

$$\int_0^{T_s} s_i(t) s_j(t) dt = 0, i \neq j$$

The minimum frequency separation between signals to make them orthogonal

is 
$$\Delta f = \frac{1}{2T_s} = \frac{R_s}{2}$$

All signals have the same energy

$$E_s = E = \int_0^{T_s} s_m(t)^2 dt$$

As a result of this condition, there will be M basis functions

#### **M-ary Coherent Frequency-Shift Keying: Signal Space Representation**

M-ary orthogonal FSK has a geometric presenation as M-dim orthogonal vectors, given as



Signals are orthogonal



#### Minimum-Distance Receiver of M-FSK



# • Modulation • $1^{*} \rightarrow s_{1}(t) = \sqrt{\frac{2E_{b}}{T_{b}}} \cos(2\pi f_{1}t)$ • $0^{*} \rightarrow s_{2}(t) = \sqrt{\frac{2E_{b}}{T_{b}}} \cos(2\pi f_{2}t)$

E<sub>b</sub> : transmitted signal energy per bit

$$\int_0^{T_b} s_1^2(t) dt = \int_0^{T_b} s_2^2(t) dt = E_b$$

- $f_i$ : transmitted frequency with separation  $\Delta f = f_1 f_0$   $\Delta f = \frac{1}{2T_h} = \frac{R_b}{2}$
- $\Delta f$  is selected so that  $s_1(t)$  and  $s_2(t)$  are orthogonal i.e.

$$\int_0^{T_b} s_1(t) s_2(t) dt = 0$$

# Example 1: Binary FSK

1.1

,

#### Two orthogonal basis functions are required

$$\phi_{1}(t) = \sqrt{\frac{2}{T_{b}}} \cos(2\pi f_{1}t) \qquad 0 \le t < T_{b} \qquad s_{1}(t) = \sqrt{E_{b}}\phi_{1}(t)$$
  
$$\phi_{2}(t) = \sqrt{\frac{2}{T_{b}}} \cos(2\pi f_{2}t) \qquad 0 \le t < T_{b} \qquad s_{2}(t) = \sqrt{E_{b}}\phi_{2}(t)$$

#### Signal space representation

$$s_{1} = \begin{bmatrix} \sqrt{E_{b}} & 0 \end{bmatrix} \qquad \qquad \underset{point \ s_{2}}{\text{Message point } s_{2}} \qquad \qquad \underset{\sqrt{E_{b}}}{\text{Message point } s_{1}} \qquad \qquad \underset$$

# Example 1: Binary FSK

#### **Observation vector**



The receiver decides in favor of  $s_1$  if the observation vector r falls inside region  $R_1$ . This occurs when  $r_1 > r_2$ 

When  $r_1 < r_2$ , *r* falls inside region R<sub>2</sub> and the receiver decides in favor of s<sub>2</sub>

# Example 1: Binary FSK

To calculate the error probability, we use the formula:

$$P_{s} \approx (\text{Number of Signals at distance dmin}) Q\left(\frac{d_{\min}}{\sqrt{2N_{0}}}\right) = (1) \cdot Q\left(\frac{\sqrt{2E_{b}}}{\sqrt{2N_{0}}}\right) = Q\left(\sqrt{\frac{E_{b}}{N_{0}}}\right)$$

where





# Example 2: 3-ary FSK

2

To calculate the error probability, we use the formula:

P<sub>s</sub> ≈ (Number of Signals at distance dmin)  $Q\left(\frac{d_{\min}}{\sqrt{2N_0}}\right) = (2) \cdot Q\left(\frac{\sqrt{2E_s}}{\sqrt{2N_0}}\right) = 2Q\left(\sqrt{\frac{E_s}{N_0}}\right)$ where

$$d_{\min} = \sqrt{2E_s}$$

#### **Minimum Distance Rule**

Calculate: 
$$(d_1)^2 = (r_1 - \sqrt{E})^2 + (r_2)^2 + (r_3)^2$$
  
 $(d_2)^2 = (r_1)^2 + (r_2 - \sqrt{E})^2 + (r_3)^2$   
 $(d_3)^2 = (r_1)^2 + (r_2)^2 + (r_3 - \sqrt{E})^2$   
Choose s1 when  $(d_1)^2 < (d_2)^2$  and  $(d_1)^2 < (d_3)$   
Equivalently, Decide s1 when

$$r_1 > r_2$$
 and  $r_1 > r_3$ 



# Error Probability in an M-ary FSK

To calculate the error probability, we use the formula:

 $P_{s} \approx (\text{Number of Signals at distance dmin}) Q\left(\frac{d_{\min}}{\sqrt{2N_{o}}}\right)$ 

$$= (\mathbf{M}-1) \cdot Q\left(\frac{\sqrt{2E_s}}{\sqrt{2N_0}}\right) = (M-1)Q\left(\sqrt{\frac{E_s}{N_0}}\right)$$

where

$$d_{\min} = \sqrt{2E_s}$$
$$E_s = (\log_2 M)E_b$$



# Bandwidth Requirements of M-FSK

• Let  $M = 2^{\lambda}$  and let the M signals be orthogonal. The minimum frequency separation between adjacent signals  $\Delta f = \frac{R_s}{2}$ .

• The bandwidth 
$$B.W = (M-1)\left(\frac{R_s}{2}\right) + 2R_s$$
.

- For the case when M = 2,  $B \cdot W = \left(\frac{R_s}{2}\right) + 2R_s = \frac{5}{2}R_s = \frac{5}{2}R_b$ .
- For the case when M = 4,  $B \cdot W = \left(\frac{3R_s}{2}\right) + 2R_s = \frac{7}{2}R_s = \frac{7}{2}\frac{R_b}{\log(4)} = \frac{7}{4}R_b$ .



#### M-ary Quadrature Amplitude Modulation (M-QAM)

 M-QAM are two-dim constellations and they involve inphase (I) and quadrature (Q) carriers:

$$\phi_I(t) = \sqrt{\frac{2}{T_s}} \cos(2\pi f_c t), \quad 0 \le t \le T_s,$$
  
$$\phi_Q(t) = \sqrt{\frac{2}{T_s}} \sin(2\pi f_c t), \quad 0 \le t \le T_s,$$

• The *i*th transmitted *M*-QAM signal is:

$$s_{i}(t) = V_{I,i} \sqrt{\frac{2}{T_{s}}} \cos(2\pi f_{c} t) + V_{Q,i} \sqrt{\frac{2}{T_{s}}} \sin(2\pi f_{c} t), \quad \substack{0 \le t \le T_{s} \\ i = 1, 2, \dots, M}$$
$$= \sqrt{E_{i}} \sqrt{\frac{2}{T_{s}}} \cos(2\pi f_{c} t - \theta_{i})$$

- $V_{I,i}$  and  $V_{Q,i}$  are the information-bearing discrete amplitudes of the two quadrature carriers,  $E_i = V_{I,i}^2 + V_{Q,i}^2$  and  $\theta_i = \tan^{-1}(V_{Q,i}/V_{I,i})$ .
- In general, QAM symbols have different energies. The average symbol energy is calculated as:

$$E_s = \sum_{i=1}^{M} E_i P[s_i(t)] = \frac{\sum_{i=1}^{M} E_i}{M}, \quad \text{for equally-likely signals}$$

- In M-QAM, the messages are encoded into both the amplitude and phase of the carrier.
- QAM is a two-dimensional encoding scheme and requires two basis functions.
  - The QAM scheme represents bits as points in a quadrant grid know as a constellation map.

 $s_i(t) = a_i \phi_1 + b_i \phi_2$  $E_i = a_i^2 + b_i^2$  (prove)

#### Criteria for Selecting a Given Constellation

- Probability of Error: In signaling over AWGN, the most likely errors are those which confuse a signal with its neighbors. To maintain the same symbol error probability, the distance between the nearest neighbors are kept the same.
- Average Transmitted Energy: The most efficient signal constellation is the one that has the smallest average transmitted energy.



**Rectangular M-QAM: Modulation and Demodulation** 



#### M-ary QAM Transmitter

- Each group of  $\lambda = \log_2 M$  bits can be divided into  $\lambda_I$  inphase bits and  $\lambda_Q$  quadrature bits, where  $\lambda_I + \lambda_Q = \lambda$ .
- Inphase bits and quadrature bits modulate the inphase and quadrature carriers *independently*.  $\{V_{ij} = -7\frac{\Delta}{2}, -5\frac{\Delta}{2}, -3\frac{\Delta}{2}, -3\frac{\Delta}{2}, -\frac{\Delta}{2}, \frac{\Delta}{2}, 3\frac{\Delta}{2}, 5\frac{\Delta}{2}, 7\frac{\Delta}{2}, 2\frac{\Delta}{2}, \frac{\Delta}{2}, \frac{\Delta}{2}$



#### M-ary QAM Receiver

Due to the orthogonality of the inphase and quadrature signals, inphase and quadrature bits can be *independently* detected at the receiver.



The most practical rectangular QAM constellation is one which  $\lambda_I = \lambda_Q = \lambda/2$ , i.e., M is a perfect square and the rectangle is a square.

#### Implementation of Rectangular M-QAM





Rectangular QAM  $M = 2^{\lambda} = 32$ Each group of  $\lambda = log_2 M$ can be divided into  $\lambda_I$  inphase bits and  $\lambda_Q$  quadrature bits where  $\lambda = \lambda_I + \lambda_Q$ . In-phase and quadrature bits modulate the in-phase and quadrature carriers independently.

#### M-QAM Constellations: Average Energy and Minimum Distance



#### M-QAM Constellations: Partitioning of the observation space



